# Computational Linguistics II: Parsing

**Bottom-up Parsing** 

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# **Example: Bottom-Up Parsing**

#### Grammar:

```
\rightarrow NP VP | NP VP kon VP
S
NP \rightarrow n
VP \rightarrow v \mid v PP \mid adv v \mid adv v PP
PP \rightarrow p NP
n → Beethoven | 1827 | this
v → expired | died
  \rightarrow in | for
kon \rightarrow and
adv \rightarrow later
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Sentence: Beethoven expired in 1827 and later died for this.

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shift: "recognize" the next word in the input

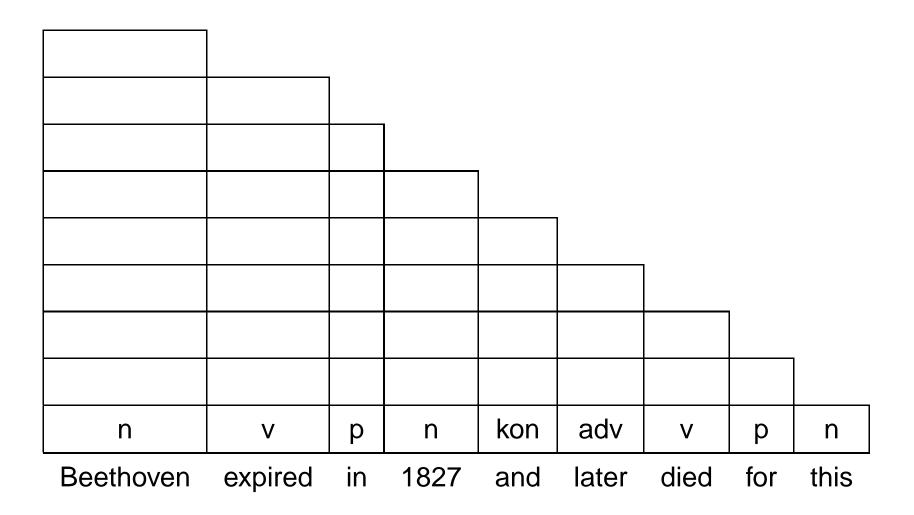
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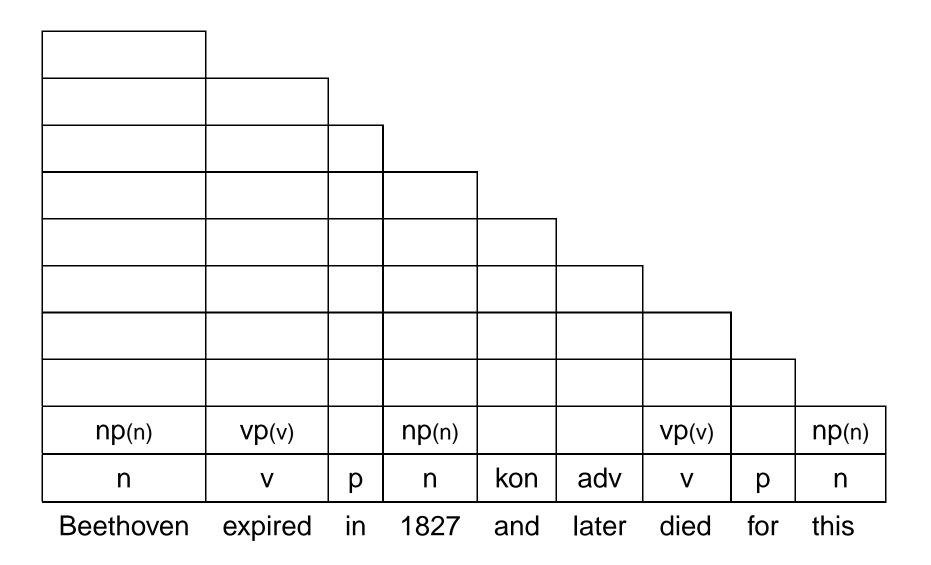
- shift: "recognize" the next word in the input
- reduce: match already recognized symbols with a righthand side of a rule and replace them by the lefthand side

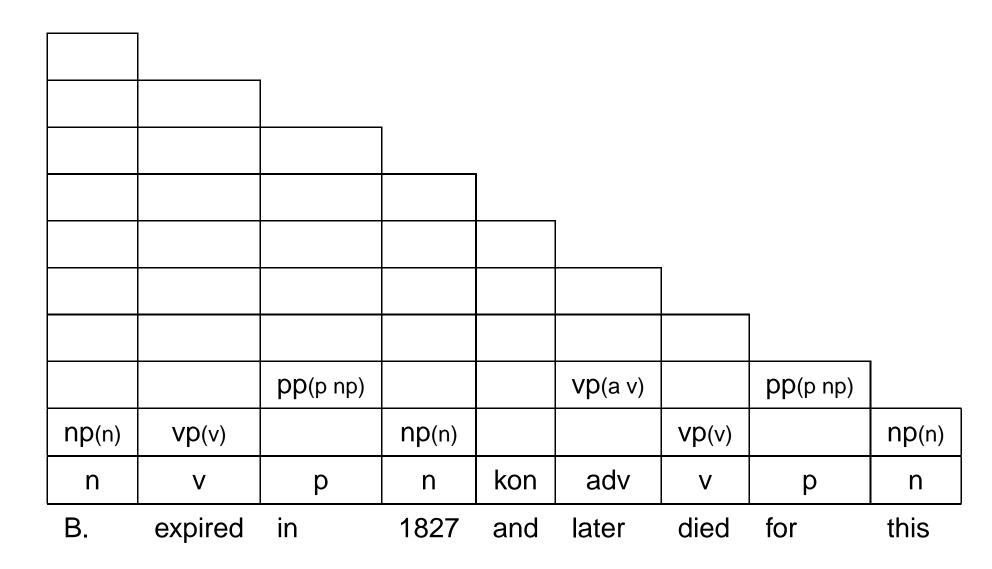
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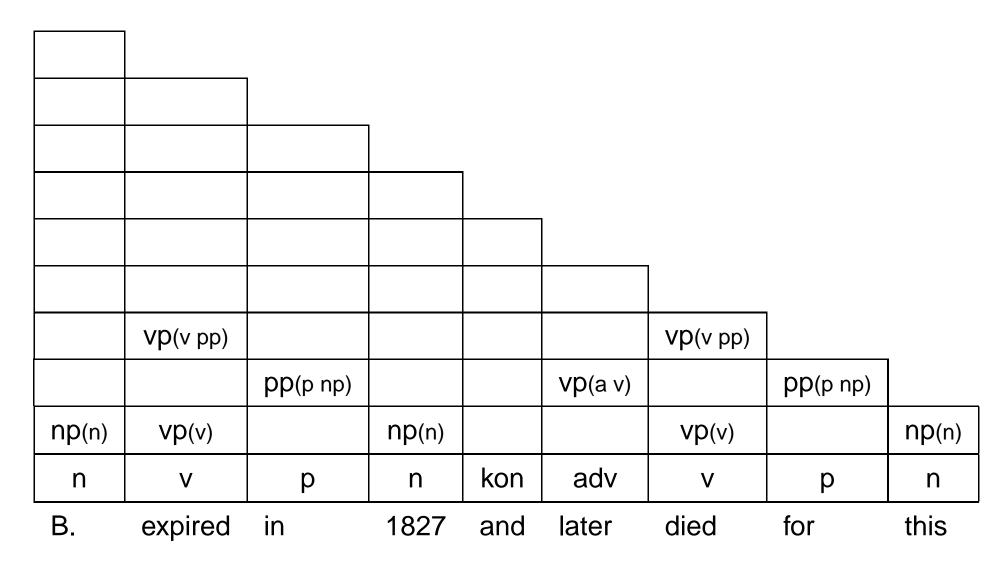
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problem: often reduces in wrong situations









	_							
S(np vp)					Vp(a v pp)			
	VP(v pp)					Vp(v pp)		
		pp(p np)			vp(a v)		pp(p np)	
np(n)	Vp(v)		np(n)			Vp(v)		np(n)
n	V	р	n	kon	adv	V	р	n
В.	expired	in	1827	and	later	died	for	this

	_							
S(np vp								
kon vp)		_						
0(1010 1110)					\/D/= = n	1		
S(np vp)					VP(a v pp)		_	
	VP(v pp)					VP(v pp)		
		pp(p np)			Vp(a v)		pp(p np)	
np(n)	Vp(v)		np(n)			Vp(v)		np(n)
n	V	р	n	kon	adv	V	р	n
В.	expired	in	1827	and	later	died	for	this

unnecessary reductions: 5 / 13 = 38.5%

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S(np vp								
kon vp)								
				]				
						1		
S(np vp)					Vp(a v pp)			
	Vp(v pp)					Vp(v pp)		
		pp(p np)			vp(a v)		pp(p np)	
np(n)	vp(v)		np(n)			vp(v)		np(n)
n	V	р	n	kon	adv	V	р	n
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problem with bottom-up: many reductions cannot be derived from start symbol e.g. VP → V NP PP the English put tacks in their tea (which started the Revolutionary Wars) no sense in grouping [tacks in their tea]

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- only VPs should be checked next!