

# Semantics 1

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# The copula verb *be*

## Different uses of *be*

- (1) Tully is Cicero.  $\rightsquigarrow$  predicative is proper noun
- (2) Cicero is a politician.  $\rightsquigarrow$  predicative is indefinite NP
- (3) Cicero is in Rome.  $\rightsquigarrow$  predicative is PP
- (4) Cicero is old.  $\rightsquigarrow$  predicative is AP

# Equative *be*

(1) Tully is Cicero.

$\text{is} \rightsquigarrow \lambda y \lambda x \lambda s . x = y$

$S$

$\lambda s . T' = C'$

NP

$T'$

N

$T'$

Tully

VP

$\lambda x \lambda s . x = C'$

V

$\lambda y \lambda x \lambda s . x = y$

is

NP

$C'$

N

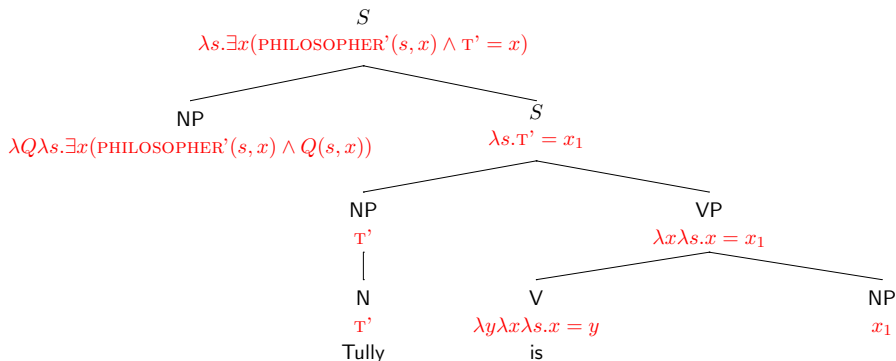
$C'$

Cicero

## Equative *be*

Equative *be* also accounts for quantifiers in predicative position.

(1) Tully is a philosopher.

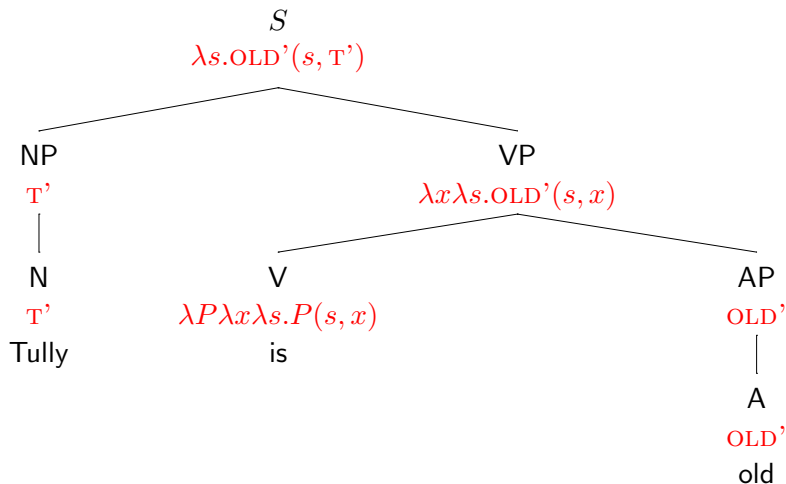


$$\lambda s. \exists x (\text{PHILOSOPHER}'(s, x) \wedge T' = x) \equiv \lambda s. \text{PHILOSOPHER}'(s, T')$$

# Predicative *be*

(1) Tully is old.

$\text{is} \rightsquigarrow \lambda P \lambda x \lambda s. P(s, x)$



# Predicative and attributive use of adjectives

- predicative use:
  - (1) Tully is old.  $\rightsquigarrow \lambda s. \text{OLD}'(s, T')$
- attributive use:
  - (2) old man  $\rightsquigarrow \lambda x \lambda s. \text{MAN}'(s, x) \wedge \text{OLD}'(s, x)$
- attributive use involves logical conjunction  $\wedge$  that is missing in predicative use
- Where does this semantic content come from?

# The syntactic solution

- Syntax:  $NP_1 \rightarrow AP, NP_2$
- Semantics:  $\|NP_1\| = \lambda x \lambda s. \|NP_2\|(s, x) \wedge \|AP\|(s, x)$
- Disadvantage:
  - does not work for all attributive adjectives:
    - (1) fake doctor
    - (2) alleged winner
    - (3) imaginary singers

# The lexical solution

## Lexical rule

If the lexicon contains an adjective  $A$  with the meaning

$$\lambda \vec{y} \lambda x \lambda s. \alpha(s, x)$$

for some predicate  $\alpha$ , then the lexicon also contains an adjective  $A$  with the meaning

$$\lambda \vec{y} \lambda P \lambda x \lambda s. P(s, x) \wedge \alpha(s, x, \vec{y})$$

**There is no consensus which solution is correct. In this course we will work with the lexical solution.**

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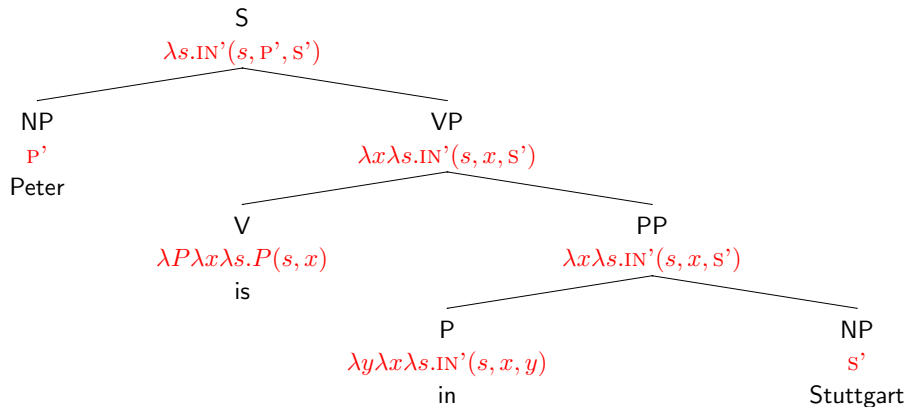
<sup>0</sup>(The notation  $\vec{y}$  represents a (possibly empty) sequence of additional arguments.)



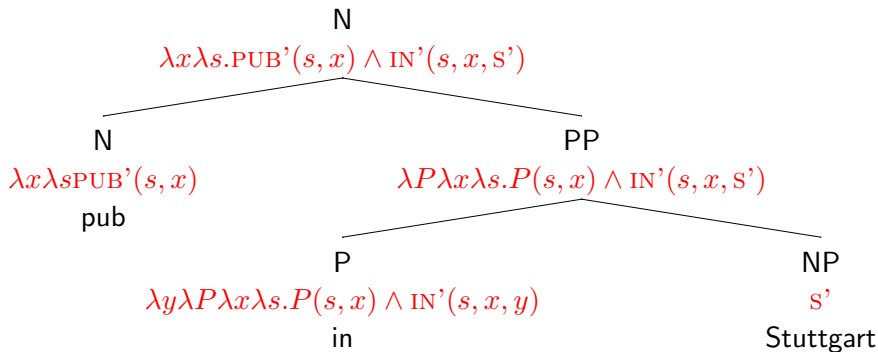
# Prepositions

- Just like APs, PPs have a predicative and a attributive use (plus an adverbial use, that will not be covered here)
- same systematic relationship between predicative and attributive use as above:
  - $\text{in}_{\text{pred}} \rightsquigarrow \lambda y \lambda x \lambda s. \text{IN}'(s, x, y)$
  - $\text{in}_{\text{attr}} \rightsquigarrow \lambda y \lambda P \lambda x \lambda s. P(s, x) \wedge \text{IN}'(s, x, y)$

# Predicative use



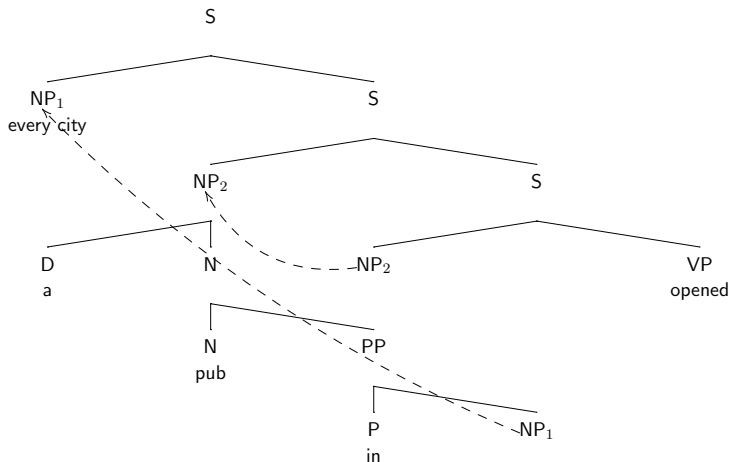
# Attributive use



# Inverse linking

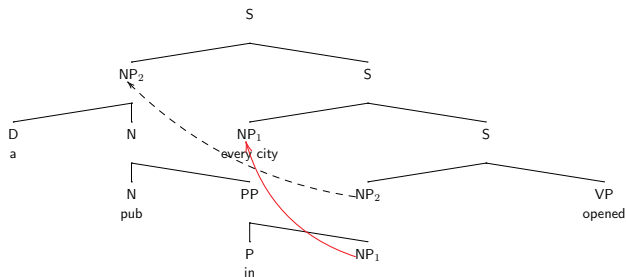
(1) A pub in every city opened.  $\rightsquigarrow$

$\lambda s. \forall y(\text{CITY}'(s, y) \rightarrow \exists x(\text{PUB}'(s, x) \wedge \text{IN}'(s, x, y) \wedge \text{OPEN}'(s, x)))$



# Inverse linking

- if we do QR in the reverse order...

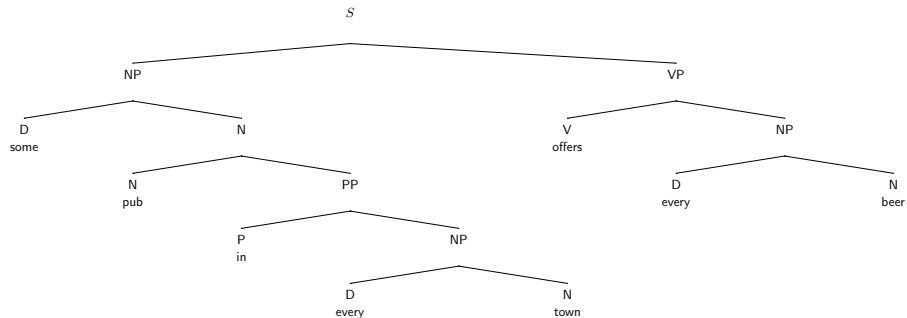


- $NP_1$  (every city) ends up not c-commanding its trace  $\Rightarrow$  illicit movement!
- semantics would come out as  $\lambda s. \exists x (\text{PUB}'(s, x) \wedge \text{IN}'(s, x, x_1) \wedge \forall y (\text{CITY}'(s, y) \rightarrow \text{OPEN}'(s, x)))$
- unbound variable (corresponds to non-c-commanded trace)

# Inverse Linking

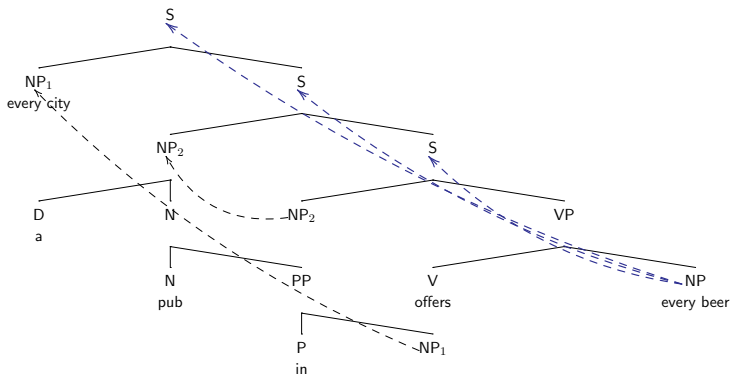
(1) Some pub in every town offers every beer.

## S-Structure



# Inverse linking

## LF 1/2/3



# Inverse Linking

(1) Some pub in every town offers every beer.

• with our current tools, we can derive three readings:

- $\lambda s. \forall z(\text{BEER}'(s, z) \rightarrow \forall y(\text{TOWN}'(s, y) \rightarrow \exists x(\text{PUB}'(s, x) \wedge \text{IN}'(s, x, y) \wedge \text{OFFER}'(s, x, z))))$
- $\lambda s. \forall y(\text{TOWN}'(s, y) \rightarrow \forall z(\text{BEER}'(s, z) \rightarrow \exists x(\text{PUB}'(s, x) \wedge \text{IN}'(s, x, y) \wedge \text{OFFER}'(s, x, z))))$
- $\lambda s. \forall y(\text{TOWN}'(s, y) \rightarrow \exists x(\text{PUB}'(s, x) \wedge \text{IN}'(s, x, y) \wedge \forall z(\text{BEER}'(s, z) \rightarrow \text{OFFER}'(s, x, z))))$

• two more readings are possible but cannot be derived so far:

- $\lambda s. \forall z(\text{BEER}'(s, z) \rightarrow \exists x(\text{PUB}'(s, x) \wedge \forall y(\text{TOWN}'(s, y) \rightarrow \text{IN}'(s, x, y)) \wedge \text{OFFER}'(s, x, z)))$
- $\lambda s. \exists x(\text{PUB}'(s, x) \wedge \forall y(\text{TOWN}'(s, y) \rightarrow \text{IN}'(s, x, y)) \wedge \forall z(\text{BEER}'(s, z) \rightarrow \text{OFFER}'(s, x, z)))$