Towards more efficient parsers

- Combining bottom-up parsing with top-down prediction
  - From shift-reduce to left-corner parsing
  - Adding more top-down filtering: link tables
- Memoization of partial results
  - well-formed substring tables
  - active charts

From shift-reduce to left-corner parsing

- Shift-reduce parsing is not goal directed at all:
  - Reduction of every possible substring,
  - obtaining every possible analysis for it.
- Idea to revise shift-reduce strategy:
  - Take a particular element x (here: the leftmost).
  - x triggers those rules it can occur in, to make predictions about the material occurring around x.

A shift-reduce parser for grammars in CNF

```
% ?- recognise([],<list(word)>,[]).

recognise([s],[],[]).

recognise([Y,X|Rest]) -->        % reduce
    {LHS ---> X,Y},
    recognise([LHS|Rest]).

recognise(Stack) -->             % shift
    [Word],
    {Cat ---> [Word]},
    recognise([Cat|Stack]).
```


A left-corner parser for grammars in CNF

```
% ?- recognise(s,<list(word)>,[]).

recognise(Phrase) --> [Word],
                      {Cat ---> [Word]},
                      lc(Cat,Phrase).

lc(Phrase,Phrase) --> [].

lc(SubPhrase,SuperPhrase) -->
    {Phrase ---> SubPhrase,Right},
    recognise(Right),
    lc(Phrase,SuperPhrase).
```

Problems of basic left-corner approach

- There can be a choice involved in picking a rule which
  - projects a particular word
  - projects a particular phrase
- How do we make sure we only pick a category which is on our path up to the goal?
  - Define a **link table** encoding the transitive closure of the left-corner relation. This is always a finite table!
  - Use it as an **oracle** guiding us to pick a reasonable candidate.

Example for a link table

For a grammar with the following non-terminal rules

```
s  ---> np, vp.        vp ---> v, np.
np ---> det, n.        n  ---> n, pp.
pp ---> p, np.
```

one can define or automatically deduce the link table

```
link(s,s).      link(np,np).    link(det,det).
link(n,n).      link(pp,pp).    link(p,p).
link(n0,n0).    link(np,s).     link(det,np).
link(p,pp).     link(v,vp).
```


Using a link table in a left-corner parser

```
recognise(Phrase) --> [Word],
                      {Cat ---> [Word]},
                      {link(Cat,Phrase)},
                      lc(Cat,Phrase).

lc(Phrase,Phrase) --> [].

lc(SubPhrase,SuperPhrase) -->
    {Phrase ---> SubPhrase,Right},
    {link(Phrase,SuperPhrase)},
    recognise(Right),
    lc(Phrase,SuperPhrase).
```


Observation: Inefficiency of backtracking

Two example sentences:

(1) He [gave [the young cat] [to Bill]].
(2) He [gave [the young cat] [some milk]].

The corresponding grammar rules:

```
vp --> v_ditrans, np, pp_to.
vp --> v_ditrans, np, np.
```

Solution: Memoization

- Store intermediate results:
  a) completely analyzed constituents:
     well-formed substring table or **(passive) chart**
  b) complete or partial analyses:
     (active) chart
- All intermediate results need to be stored for completeness.
- All possible solutions are explored in parallel.

CYK Parser

- Developed independently by Cocke, Younger, and Kasami
- Grammar has to be in Chomsky Normal Form (CNF), only
  - RHS with a single terminal: $A \rightarrow a$
  - RHS with two non-terminals: $A \rightarrow BC$
- The well-formed substring table, henceforth (passive) chart, for a string of length n is an $n \times n$ matrix.

The passive chart

- Sentence representation showing position and word indices:

$$\cdot_0 \; \mathrm{w}_1 \; \cdot_1 \; \mathrm{w}_2 \; \cdot_2 \; \mathrm{w}_3 \; \cdot_3 \; \mathrm{w}_4 \; \cdot_4 \; \mathrm{w}_5 \; \cdot_5 \; \mathrm{w}_6 \; \cdot_6$$

- An entry in a field (i, j) of the chart encodes the set of categories which spans the string from position i to j.
- More formally: chart(i,j) $= \{A \mid A \Rightarrow^* w_{i+1} \ldots w_j\}$

Coverage represented in the chart

An input sentence with 6 words:

$$\cdot_0 \; \mathrm{w}_1 \; \cdot_1 \; \mathrm{w}_2 \; \cdot_2 \; \mathrm{w}_3 \; \cdot_3 \; \mathrm{w}_4 \; \cdot_4 \; \mathrm{w}_5 \; \cdot_5 \; \mathrm{w}_6 \; \cdot_6$$

Coverage represented in the chart:

TO:

FROM:	1	2	3	4	5	6
0	0–1	0–2	0–3	0–4	0–5	0–6
1		1–2	1–3	1–4	1–5	1–6
2			2–3	2–4	2–5	2–6
3				3–4	3–5	3–6
4					4–5	4–6
5						5–6

Example for coverage represented in chart

Example sentence:

$$\cdot_0 \text{ the } \cdot_1 \text{ young } \cdot_2 \text{ boy } \cdot_3 \text{ saw } \cdot_4 \text{ the } \cdot_5 \text{ dragon } \cdot_6$$

Coverage represented in chart:

	1	2	3	4	5	6
0	the	the young	the young boy	the young boy saw	the young boy saw the	the young boy saw the dragon
1		young	young boy	young boy saw	young boy saw the	young boy saw the dragon
2			boy	boy saw	boy saw the	boy saw the dragon
3				saw	saw the	saw the dragon
4					the	the dragon
5						dragon

An example for a filled-in chart

Input sentence:
$\cdot_0$ the $\cdot_1$ young $\cdot_2$ boy $\cdot_3$ saw $\cdot_4$ the $\cdot_5$ dragon $\cdot_6$

Chart:

	1	2	3	4	5	6
0	{Det}	{}	{NP}	{}	{}	{S}
1		{Adj}	{N}	{}	{}	{}
2			{N}	{}	{}	{}
3				{V}	{}	{VP}
4					{Det}	{NP}
5						{N}

Grammar:
S → NP VP
VP → Vt NP
NP → Det N
N → Adj N
Vt → saw
Det → the
Det → a
N → dragon
N → boy
Adj → young

Filling in the chart left-to-right, depth-first

	1	2	3	4	5	6
0	**1!**	**3**	**6**	**10**	**15**	**21**
1		**2!**	**5**	**9**	**14**	**20**
2			**4!**	**8**	**13**	**19**
3				**7!**	**12**	**18**
4					**11!**	**17**
5						**16!**

for $j := 1$ to 6
  lexical-chart-fill$(j-1, j)$
  for $i := j-2$ down to 0
    syntactic-chart-fill(i, j)

lexical-chart-fill(j-1,j)

- Idea: Lexical lookup. Fill the field $(j-1, j)$ in the chart with the preterminal category dominating word j.
- Realized as:

$$chart(j-1, j) := \{\mathsf{X} \mid \mathsf{X} \rightarrow \mathsf{word}_j \in \mathsf{P}\}$$

syntactic-chart-fill(i,j)

- Idea: Perform all reduction step using syntactic rules such that the reduced symbol covers the string from i to j.
- Realized as:

$$chart(i,j) = \left\{ A \;\middle|\; \begin{array}{l} A \rightarrow BC \in P, \\ i < k < j, \\ B \in chart(i,k), \\ C \in chart(k,j) \end{array} \right\}$$

Explicit version of syntactic-chart-fill(i,j)

- Needed: version making explicit enumerations of
  - every possible value of k and
  - every context free rule
- Code:
  $chart(i,j) := \{\}$.
  for $k := i+1$ to $j-1$ do
  for every $A \rightarrow BC \in P$ do
  if $B \in chart(i,k)$ and $C \in chart(k,j)$ then
  $chart(i,j) := chart(i,j) \cup \{\text{A}\}$.

The complete CYK algorithm

for $j := 1$ to n do
    $chart(j-1,j) := \{\text{X} \mid \text{X} \rightarrow \text{word}_j \in \text{P}\}$
    for $i := j-2$ down to 0 do
        $chart(i,j) := \{\}$
        for $k := i+1$ to $j-1$ do
            for every $A \rightarrow BC \in P$ do
                if $B \in chart(i,k)$ and $C \in chart(k,j)$ then
                    $chart(i,j) := chart(i,j) \cup \{\text{A}\}$
if S $\in chart(0,n)$ then accept else reject

The CYK algorithm in PROLOG
(cky/cky.pl)

```
% Data structures: chart(From,To,Category)
:- dynamic chart/3.

% Operator for grammar rules
:- op(1200,xfx,'--->').
```

```
% recognize(+WordList,?Startsymbol)
% top-level predicate for CYK recognizer

recognize(S,Cat) :-
   retractall(chart(_,_,_)),
   length(S,N),
   fill(0,N,S),
   chart(0,N,Cat).
```



```
% fill(+Current minus one,+Last,+WordList)
% Main j-loop from 1 to number of words in string.

fill(N,N,[]).
fill(JminOne,N,[W|Ws]) :-
   J is JminOne + 1,
   lexical_chart_fill(J,JminOne,W),
   %
   I is J - 2,
   syntactic_chart_fill(I,J),
   %
   fill(J,N,Ws).
```



```
% lexical_chart_fill(+J,+JminOne,+Word)
% fill main diagonal with preterminal categories

lexical_chart_fill(J,JminOne,W) :-
   findall_unique(X,(X ---> [W]),Cats),
   add_all_to_chart(JminOne,J,Cats).
```



```
% syntactic_chart_fill(+I,+J)
% i-loop from J-2 down to 0

syntactic_chart_fill(-1,_) :- !.
syntactic_chart_fill(I,J) :-
   K is I+1,
   build_phrases_from_to(I,K,J),
   IminOne is I-1,
   syntactic_chart_fill(IminOne,J).
```

```
% build_phrases_from_to(+From,+Current,+To)

build_phrases_from_to(_,J,J) :- !.
build_phrases_from_to(I,K,J) :-
   findall_unique(A,(chart(I,K,B),
                     chart(K,J,C),
                     (A ---> [B,C])),
                  List),
   add_all_to_chart(I,J,List),
   KplusOne is K+1,
   build_phrases_from_to(I,KplusOne,J).
```



```
% add_one_to_chart(+FromIndex,+ToIndex,+Contents)
% a) only add if it does not yet exist:
add_one_to_chart(From,To,Cat) :- chart(From,To,Cat), !.

% b) add a chart entry
add_one_to_chart(From,To,Cat) :-
   assertz(chart(From,To,Cat)).

add_all_to_chart(_,_,[]).
add_all_to_chart(From,To,[Cat|Cats]) :-
   add_one_to_chart(From,To,Cat),
   add_all_to_chart(From,To,Cats).
```